

Class forcing in its rightful setting

Victoria Gitman

vgitman@nylogic.org
<http://victoriagitman.github.io>

KGRC Seminar

June 25, 2020

Class forcing in set theory

Class forcing (forcing with a class partial order) is **ubiquitous** in set theory.

- Make **GCH fail at every regular cardinal**.
 - ▶ $\prod_{\alpha \in \text{Reg}} \text{Add}(\alpha, \alpha^{++})$ with Easton support.
- Make $V = \text{HOD}$ hold by **coding all sets into the continuum function**.
- Make all **supercompact cardinals Laver indestructible**.
- Add a **global well-order** (bijection $F : V \xrightarrow[\text{onto}]{1-1} \text{Ord}$).
 - ▶ **Add(Ord, 1)**: conditions are partial functions $f : \alpha \rightarrow \text{Ord}$ for $\alpha \in \text{Ord}$ ordered by extension.
 - ▶ **<Ord-closed**.
 - ▶ Does **not** add sets.
- **Shoot a class club through a fat stationary class**.
 - ▶ Conditions are closed sets of ordinals from the stationary class ordered by end-extension.
 - ▶ **<Ord-distributive**.
 - ▶ Does **not** add sets.

The right setting for class forcing

Class forcing is **fundamentally about classes**.

- A **generic filter** for a class partial order is a **subclass** of the partial order meeting all dense **subclasses**.
- The properties of the partial order **depend on which classes exist around it**.

First-order set theory

- Classes are **definable** (with parameters) collections of sets.
- Classes are objects in the **meta-theory**.

Second-order set theory

- Classes are **elements of the model**.
- We can **quantify** over classes.
- We can **study general properties** of classes.
- **The theory determines which classes exist**.

We should not expect class forcing to be as nice as set forcing because classes even in the strongest second-order set theories do not behave as nicely as sets.

Second-order set theory

Second-order set theory has two sorts of objects: **sets** and **classes**.

Syntax: Two-sorted logic

- Separate variables and quantifiers for sets and classes
- Convention: upper-case letters for classes, lower-case letters for sets
- Notation:
 - ▶ Σ_n^0 - first-order Σ_n -formula
 - ▶ Σ_n^1 - n -alternations of class quantifiers followed by a first-order formula

Semantics: A model is a triple $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle$.

- V consists of the **sets**.
- \mathcal{C} consists of the **classes**.
- Every set is a class: $V \subseteq \mathcal{C}$.
- $C \subseteq V$ for every $C \in \mathcal{C}$.

Weakest second-order theories

Set axioms

- ZF^- : no choice, no powerset, collection instead of replacement
- ZFC^- : no powerset, collection instead of replacement, well-ordering principle instead of axiom of choice¹

¹Different formulations of the axiom of choice are not equivalent over ZF^- .

Weakest second-order set theories

Second-order theories: axioms for sets and classes

- **GB**
 - ▶ set axioms: ZF
 - ▶ **Replacement**: If F is a function and a is a set, then $F \upharpoonright a$ is a set.
 - ▶ **First-order comprehension**: If $\varphi(x, A)$ is a first-order formula, then $\{x \mid \varphi(x, A)\}$ is a class.
 - ▶ Every model of ZF with together with its definable collections is a model of GB.
- **GBc**: GB with ZF replaced by ZFC.
 - ▶ Every model of ZFC together with its definable collections is a model of GBc.
- **GBC**: Gödel-Bernays set theory
 - ▶ GBc
 - ▶ **Global well-order**: there is a bijection $F : V \xrightarrow[\text{onto}]{1-1} \text{Ord}$.²
 - ▶ Every model of ZFC, with a definable global well-order, together with its definable collections is a model of GBC, e.g. L .
 - ▶ Every model of ZFC has a class forcing extension with the same sets satisfying GBC: force with $\text{Add}(\text{Ord}, 1)$.
 - ▶ **Conservative** over ZFC.
- **GB⁻**, **GBc⁻**, **GBC⁻**: ZF⁻ instead of ZF.

²Different formulations of global choice are not equivalent over GBC^- .

Hyperclasses

Definition: In a model of second-order set theory, a **hyperclass** is a **definable** (with parameters) collection of classes.

- Analogue of classes in second-order.
- Third-order part of the model.

Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$. Given a set a and a class A , the **a -th slice of A** is the class

$$\{x \mid \langle a, x \rangle \in A\},$$

where $\langle \cdot, \cdot \rangle$ is the Gödel pairing function.

Definition: A **hyperclass**, given by a formula $\varphi(X, A)$, is **coded by a class C** if for every class B , $\varphi(B, A)$ holds if and only if **there is a set b such that $C_b = B$** .

“A coded hyperclass has class-many classes.”

Transfinite recursion on classes along meta-ordinals

Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$.

Definition: A **meta-ordinal** is a **well-order** $(\Gamma, \leq) \in \mathcal{C}$.

- Examples: Ord , $\text{Ord} + \text{Ord}$, $\text{Ord} \cdot \omega$.
- Notation: For $a \in \Gamma$, $\Gamma \upharpoonright a$ is the restriction of the well-order to \leq -predecessors of a .
- Meta-ordinals may not have unique representations.

Question: Does **GBC** prove that **any two meta-ordinals are comparable**?

Definition: Suppose $\Gamma \in \mathcal{C}$ is a **meta-ordinal**. A **solution along Γ** to a **first-order recursion rule** $\varphi(x, b, F)$ is a class S such that for every $b \in \Gamma$, $S_b = \{x \mid \varphi(x, b, S \upharpoonright b)\}$.

- $S_0 = \{x \mid \varphi(x, 0, \emptyset)\}$
- $S_1 = \{x \mid \varphi(x, 1, \langle S_0 \rangle)\}$
- $S_2 = \{x \mid \varphi(x, 2, \langle S_0, S_1 \rangle)\}$
- etc.

ETR: Elementary transfinite recursion

ETR: Every first-order recursion on classes along a meta-ordinal has a solution.

ETR $_{\Gamma}$: ETR restricted to a well-order Γ .

- ETR $_{\text{Ord}\cdot\omega}$, ETR $_{\text{Ord}}$, ETR $_{\omega}$.

Theorem: (Fujimoto) Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$.

- ETR $_{\omega}$ is equivalent to the assertion that for every class A , there is a truth predicate for the structure $\langle V, \in, A \rangle$ (with a predicate for A).
 - ▶ Tarskian truth is given by a recursion of length ω each of whose levels is a class.
- ETR $_{\Gamma}$ is equivalent to the existence, for every class A , of the iterated truth predicate of length Γ for the structure $\langle V, \in, A \rangle$.

Corollary: ETR is equivalent to the assertion that for every meta-ordinal Γ and every class A , there is an iterated truth predicate of length Γ for $\langle V, \in, A \rangle$.

Corollary: GBC + ETR $_{\omega}$ implies Con(ZFC).

Theorem: (Williams) If $\Gamma > \omega^{\omega}$ is a meta-ordinal, then GBC + ETR $_{\Gamma\cdot\omega}$ implies Con(GBC + ETR $_{\Gamma}$).

Theorem: GBC + ETR implies that any two meta-ordinals are comparable.

A comprehension hierarchy to Kelley-Morse set theory

Σ_n^1 -**comprehension CA**: Every Σ_n^1 -formula defines a class.

Theorem: $\text{GBC} + \Sigma_1^1\text{-CA}$ implies ETR .

The theories $\text{GBC} + \Sigma_n^1\text{-CA}$ form a hierarchy of strength culminating in Kelley-Morse set theory KM .

Kelley-Morse set theory KM

- GBC
- Every second-order formula defines a class.

Class choice principles

Choice Scheme CC: Given a second-order formula $\varphi(x, X, A)$, if for every set x , there is a class X witnessing $\varphi(x, X, A)$, then there is a class collecting witnesses for every x :

$$\forall x \exists X \varphi(x, X, A) \rightarrow \exists Y \forall x \varphi(x, Y_x, A).$$

Σ_n^1 -CC: CC restricted to Σ_n^1 -formulas.


Theorem: (G., Hamkins) Σ_1^0 -CC can fail in a model of KM.

α -Dependent Choice Scheme DC_α : ($\omega \leq \alpha \leq \text{Ord}$) Every second-order definable $<\alpha$ -closed tree $T(X, A)$ has a branch of height α .

Σ_n^1 - DC_α : DC_α restricted to Σ_n^1 -definable trees.

Theorem: (Marek, Mostowski, Ratajczyk) If $\mathcal{V} \models \text{GBC} + \Sigma_n^1\text{-CA}$, then its second-order constructible universe $\mathcal{L} \models \text{GBC} + \Sigma_n^1\text{-Comprehension} + \Sigma_n^1\text{-CC} + \Sigma_n^1\text{-DC}_{\text{Ord}}$.

- If $\mathcal{V} \models \text{KM}$, then its second-order constructible universe $\mathcal{L} \models \text{KM} + \text{CC} + \text{DC}_{\text{Ord}}$.
 - ▶ Given a meta-ordinal Γ , we can build a meta-constructible universe L_Γ by a recursion of length Γ .
 - ▶ A meta-ordinal Γ is **constructible** if $\Gamma \in L_{\text{Ord}^+}$.
 - ▶ **Theorem:** (Tharp) Constructible meta-ordinals have unique representations.
 - ▶ A class $A \in \mathcal{C}$ is **constructible** if there is a constructible meta-ordinal Γ such that $A \in L_{\text{Ord}^+}$.
 - ▶ The **second-order constructible universe** is $\mathcal{L} = \langle L, \in, \mathcal{L} \rangle$, where \mathcal{L} consists of the constructible classes.

- The theories **KM** and **KM+CC+DC_{Ord}** are **equiconsistent**. 

Applications of class choice

Theorem: Over GBC, $\Sigma_n^1\text{-CC}$ implies $\Delta_n^1\text{-CA}$. Consequently, $\text{GBC}+\text{CC}$ is equivalent to $\text{KM}+\text{CC}$.

Theorem: (G., Hamkins) Over GBC, $\text{Loś' Lemma for second-order ultrapowers}$ is equivalent to $\text{CC for set-many choices}$.

Theorem: (G., Hamkins) Over GBC, $\text{normal form for second-order formulas}$ is equivalent to CC .

Theorem: (G., Hamkins, Karagila) $\text{Fodor's Lemma for class clubs}$

- can fail in KM ,
- holds in $\text{GBC}+\Sigma_2^0\text{-CC}$.

Aside

Theorem: (Enayat) $\text{GBC}+\Sigma_1^1\text{-CC}$ is conservative over ZFC .

- Every countable model of ZFC can be elementarily extended to countable recursively saturated model of ZFC .
- Every countable recursively saturated model of ZFC can be expanded to a model of $\text{GBC} + \Sigma_1^1\text{-CC}$ without adding sets.

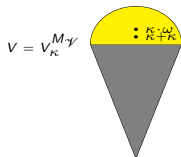
Moving to first-order with KM+CC

Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{KM+CC}$.

- View each **extensional well-founded** class relation $R \in \mathcal{C}$ as **coding a transitive set**.
 - ▶ $\text{Ord} + \text{Ord}$, $\text{Ord} \cdot \omega$
 - ▶ $V \cup \{V\}$
- Define a **membership relation E** on the collection of all such relations R (modulo isomorphism).
- Let $\langle M_{\mathcal{V}}, E \rangle$, **the companion model of \mathcal{V}** , be the resulting **first-order structure**.
 - ▶ $M_{\mathcal{V}}$ has the **largest cardinal** $\kappa \cong \text{Ord}^{\mathcal{V}}$.
 - ▶ $V_{\kappa}^{M_{\mathcal{V}}} \cong V$.
 - ▶ $\mathcal{P}(V_{\kappa})^{M_{\mathcal{V}}} \cong \mathcal{C}$.
 - ▶ $\langle M_{\mathcal{V}}, E \rangle \models \text{ZFC}_{\Gamma}^{-}$.

ZFC $_{\Gamma}^{-}$

- ZFC $^{-}$
- There is a **largest cardinal** κ .
- κ is **inaccessible**: κ is regular and for all $\alpha < \kappa$, 2^{α} exists and $2^{\alpha} < \kappa$.
 - ▶ V_{κ} exists.
 - ▶ $V_{\kappa} \models \text{ZFC}$.



Bi-interpretability

Suppose $M \models \text{ZFC}_I^-$ with a largest cardinal κ .

- $V = V_\kappa^M$
- $\mathcal{C} = \{X \in M \mid X \subseteq V_\kappa^M\}$
- $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{KM} + \text{CC}$
- $M_{\mathcal{V}} \cong M$ is the companion model of \mathcal{V} .

Theorem: (Marek) The theory $\text{KM} + \text{CC}$ is bi-interpretable with the theory ZFC_I^- .

Class forcing interpretation

Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$ and $\mathbb{P} \in \mathcal{C}$ is a **class forcing**.

Definition:

- A **class \mathbb{P} -name** is a collection of pairs $\langle \sigma, p \rangle$ such that $p \in \mathbb{P}$ and $\sigma \in V^{\mathbb{P}}$.
- G is **\mathcal{V} -generic for \mathbb{P}** if G meets every **dense subclass** $D \in \mathcal{C}$ of \mathbb{P} .
- The forcing extension $\mathcal{V}[G] = \langle V[G], \in, \mathcal{C}[G] \rangle$.

Definability of forcing relations

Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$ and $\mathbb{P} \in \mathcal{C}$ is a class forcing.

The Class Forcing Theorem for \mathbb{P} : There is a **solution** to the **recursion** defining the forcing relation for atomic formulas.

- $p \Vdash \sigma \in \tau$: there is a dense set of conditions $q \leq p$ for which there is $\langle \rho, r \rangle \in \tau$ with $q \leq r$ and $q \Vdash \sigma = \rho$.
- $p \Vdash \sigma = \tau$: $p \Vdash \sigma \subseteq \tau$ and $p \Vdash \tau \subseteq \sigma$.
- $p \Vdash \sigma \subseteq \tau$: whenever $\langle \rho, r \rangle \in \sigma$ and $q' \leq p, r$, there is $q \leq q'$ with $q \Vdash \rho \in \tau$.

The **Class Forcing Theorem** implies that **forcing relations for all second-order formulas** are **definable**.

- $p \Vdash \sigma \in \Gamma$: there are densely many $q \leq p$ for which there is $\langle \tau, r \rangle \in \Gamma$ with $q \leq r$ and $q \Vdash \sigma = \tau$.
- $p \Vdash \varphi \wedge \psi$: $p \Vdash \varphi$ and $p \Vdash \psi$.
- $p \Vdash \neg \varphi$: there is no $q \leq p$ with $q \Vdash \varphi$.
- $p \Vdash \forall x \varphi(x)$: $p \Vdash \varphi(\tau)$ for every \mathbb{P} -name τ .
- $p \Vdash \forall X \varphi(X)$: $p \Vdash \varphi(\Delta)$ for every class \mathbb{P} -name Δ .

Nice class forcing

Definition: A class forcing \mathbb{P} has the **Ord-cc** if every **antichain** of \mathbb{P} is a **set**.

Definition: (Friedman) A class forcing \mathbb{P} is **pretame** if for every class sequence $\langle D_x \mid x \in a \rangle \in \mathcal{C}$ of dense classes of \mathbb{P} , indexed by elements of a set a , and condition $p \in \mathbb{P}$, there is a condition $q \leq p$ and a sequence $\langle d_x \mid x \in a \rangle$ of subsets of \mathbb{P} such that each $d_x \subseteq D_x$ is pre-dense below q in \mathbb{P} .

- “Reduces dense classes to pre-dense sets.”
- Condition to preserve replacement

Definition: (Friedman) A class forcing \mathbb{P} is **tame** if it is pretame and for every $p \in \mathbb{P}$, there is $q \leq p$ and ordinal α such that whenever $\vec{D} = \{\langle D_0^x, D_1^x \rangle \mid x \in a\} \in \mathcal{C}$, for a set a , is a sequence of pre-dense partitions below q , then the class

$$\{r \in \mathbb{P} \mid \vec{D} \text{ is equivalent below } r \text{ to some partition } \vec{E} \in V_\alpha\}$$

is dense below q .

- Condition to preserve powerset
- Ord-cc forcings
- Progressively closed Ord-length products and iterations of set forcing
- $<$ Ord-distributive forcing

Preserving the theory

Class forcing **need not preserve ZFC**.

- $\text{Coll}(\omega, \text{Ord})$.
- (Friedman) Forcing \mathbb{F} to code $\langle V, \in \rangle$ into a relation E on ω .
 - ▶ In the extension: $\langle V, \in \rangle \cong \langle \omega, E \rangle$.

Question: Are class forcing extensions always closed under complements?

Theorem: (Friedman) **Pretameness** is equivalent to **preservation of GB^-** . **Tameness** is equivalent to **preservation of GB** .

Theorem: **Tame** forcings **preserve**:

- GBC , GBC .
- (Antos) KM .
- (Antos, G., Friedman) $\text{KM}+\text{CC}$.
- (Antos, G., Friedman) $\text{KM}+\text{CC}+\text{DC}_\alpha$ for every $\omega \leq \alpha \leq \text{Ord}$.

Theorem: **Tame** forcings **preserve $\text{GBC} + \Sigma_n^1\text{-CA} + \Sigma_n^1\text{-CC}$** .

Question: Do **tame** forcings **preserve $\text{GBC} + \Sigma_n^1\text{-CA}$** ?

- Without $\Sigma_n^1\text{-AC}$, the **forcing relation for Σ_n^1 -formulas** may not be Σ_n^1 -definable.

Class Forcing Theorem

Theorem: (Holy, Krapf, Lücke, Njegomir, Schlicht) The [Class Forcing Theorem](#) can fail in a model of GBC .

- The [Class Forcing Theorem](#) for Friedman's forcing \mathbb{F} implies existence of [truth predicate](#) for $\langle V, \in \rangle$.

Theorem: (Stanley) In GB^- , the [Class Forcing Theorem](#) holds for all pretame forcings.

- In ZFC , pretame forcings have definable forcing relations.
- In ZFC^- pretame forcings have definable forcing relations.

Theorem: (G., Hamkins, Holy, Schlicht, Williams) Over GBC , the [Class Forcing Theorem](#) is [equivalent](#) to ETR_{Ord} .

- The [recursion](#) to define the forcing relation on [atomic formulas](#) has length Ord .
- The [atomic forcing relation](#) for Friedman's forcing \mathbb{F} yields an [iterated truth predicate](#) of length Ord for $\langle V, \in \rangle$.
- Define a version of Friedman's forcing \mathbb{F}_A for a class A .

Dense embeddings

Theorem: Suppose \mathbb{P} and \mathbb{Q} are set forcings and there is a dense embedding from \mathbb{P} into \mathbb{Q} . Then \mathbb{P} and \mathbb{Q} have the same forcing extensions.

Definition:

- $\text{Coll}(\omega, \text{Ord})$: conditions $f : A \rightarrow \text{Ord}$, where A is a finite subset of ω .
- $\text{Coll}_*(\omega, \text{Ord})$: conditions $f : n \rightarrow \text{Ord}$ (no holes!).

Theorem (Holy, Krapf, Lücke, Njegomir, Schlicht)

- $\text{Coll}_*(\omega, \text{Ord})$ densely embeds into $\text{Coll}(\omega, \text{Ord})$.
- $\text{Coll}(\omega, \text{Ord})$ adds a bijection between every ordinal and ω .
- $\text{Coll}_*(\omega, \text{Ord})$
 - ▶ adds a class $f : \omega \xrightarrow[\text{onto}]{1-1} \text{Ord}$,
 - ▶ does not add sets.

Theorem: (Holy, Krapf, Lücke, Njegomir, Schlicht) Assume GBC^- . If \mathbb{P} and \mathbb{Q} are pretame class forcings such that \mathbb{P} densely embeds into \mathbb{Q} , then \mathbb{P} and \mathbb{Q} have the same forcing extensions.

Nice names

Definition: Suppose \mathbb{P} is a forcing. A **nice \mathbb{P} -name for a subset of ordinals** has the form $\bigcup_{\xi < \alpha} \{\check{\xi}\} \times A_\xi$, where each $A_\xi \subseteq \mathbb{P}$ is an **antichain**.

Theorem: Suppose \mathbb{P} is a **set forcing**. Then every **subset of ordinals** in a forcing extension by \mathbb{P} has a **nice \mathbb{P} -name**.

Theorem: (Holy, Krapf, Schlicht) Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$ and $G \subseteq \text{Coll}(\omega, \text{Ord})$ is \mathcal{V} -generic. In $\mathcal{V}[G]$, the set

$$A = \{n \in \omega \mid G(n) = 0\}$$

does not have a nice- $\text{Coll}(\omega, \text{Ord})$ -name.

Theorem: (Holy, Krapf, Schlicht)

- Assume GBC^- . Every **pretame** class forcing \mathbb{P} has **nice \mathbb{P} -names for all subsets of ordinals**.
- Assume $\text{GBC} + \text{ETR}_{\text{Ord}}$. A class forcing \mathbb{P} that has **nice \mathbb{P} -names for all subsets of ordinals** is **pretame**.
 - ▶ Uses existence of **definable forcing relation**.

Ground model definability

Theorem: (Laver, Woodin) The ground model V is uniformly definable from a parameter in any set-forcing extension.

Theorem: (Antos) Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$. There is a class forcing \mathbb{P} such that $\langle V, \in \rangle$ is not definable from parameters over $V[G]$ in any forcing extension $\mathcal{V}[G]$ by \mathbb{P} .

- \mathbb{P} is the Easton support product $\prod_{\alpha \in \text{Reg}} \text{Add}(\alpha, 1)$.
- $\text{Add}(\alpha, 1) \cong \text{Add}(\alpha, 1) \times \text{Add}(\alpha, 1)$.
- Use an automorphism.

Theorem: (G., Johnstone) It is consistent that $H_{\kappa^+} (\models \text{ZFC}_1^-)$ is not definable, even with parameters, in its forcing extension $H_{\kappa^+}[G]$ by $\text{Add}(\kappa, 1)$.

Corollary: There is a model $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{KM} + \text{CC}$ such that \mathcal{C} is not a hyperclass of its forcing extensions $\mathcal{V}[G] = \langle V[G], \in, \mathcal{C}[G] \rangle$ by $\text{Add}(\text{Ord}, 1)$, even with a parameter from $\mathcal{C}[G]$.

Theorem: (Asperó) It is consistent (from large cardinals) that $H_{\kappa^+} (\models \text{ZFC}_1^-)$ is not definable, even with parameters, in its forcing extension $H_{\kappa^+}[G]$ by $\text{Add}(\omega, 1)$.

Corollary: Ground model definability can fail in $\text{KM} + \text{CC}$ even for set-forcing.

Boolean completions

Theorem: Every set forcing \mathbb{P} densely embeds into a unique (modulo isomorphism) complete Boolean algebra.

- $U \subseteq \mathbb{P}$ is a **cut** if it is **closed downwards**: if $p \in U$ and $q \leq p$, then $q \in U$.
- For $p \in \mathbb{P}$, $U_p = \{q \in \mathbb{P} \mid q \leq p\}$ is a **cut**.
- A **cut** U is **regular** if for every $p \notin U$, there is $q \leq p$ such that $U_q \cap U = \emptyset$.
- The Boolean algebra $\mathbb{B}_{\mathbb{P}}$ consists of all **regular cuts** of \mathbb{P} .

Definition: A class Boolean algebra is:

- **set-complete** if it has **suprema** for all its **subsets**,
- **class-complete** if it has **suprema** for all its **subclasses**.

Theorem: (Holy, Krapf, Lücke, Njegomir, Schlicht) Assume GBC. A class forcing has:

- a **Boolean set-completion** if and only if the **Class Forcing Theorem holds** for it.
 - ▶ The forcing relation can be used to construct a Boolean set-completion.
 - ▶ Defining Boolean values of assertions in the forcing language does not require ETR.
- a **Boolean class-completion** if and only if it has the **Ord-cc**.
- a **unique Boolean set-completion** if and only if it has the **Ord-cc**.

A class forcing with a proper class antichain cannot have a Boolean class-completion!

Hyperclass Boolean completions

Definition: Suppose $\mathcal{V} \models \text{GBC}$.

- Suppose $\mathbb{P} \in \mathcal{C}$ is a class forcing. The regular cuts of \mathbb{P} form a **hyperclass Boolean algebra** $\mathbb{B}_{\mathbb{P}}$.
- A **hyperclass Boolean algebra** is:
 - ▶ **class-complete** if it has **suprema** for all its **coded subhyperclasses**.
 - ▶ **hyperclass-complete** if it has **suprema** for all its **subhyperclasses**.

Theorem: Suppose $\mathcal{V} \models \text{GBC}$ and $\mathbb{P} \in \mathcal{C}$ is a class forcing. Then $\mathbb{B}_{\mathbb{P}}$ is **class-complete**.

Theorem: (Antos, Friedman, G.) Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$ and $\mathbb{P} \in \mathcal{C}$ is **not Ord-cc**. If $\mathbb{B}_{\mathbb{P}}$ is **hyperclass-complete**, then $\mathcal{V} \models \text{KM}$.

- Code instances of comprehension into suprema of antichains of $\mathbb{B}_{\mathbb{P}}$.

Theorem: (Antos, Friedman, G.) Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{KM}$ and $\mathbb{P} \in \mathcal{C}$ is a class forcing. Every **hyperclass antichain** of $\mathbb{B}_{\mathbb{P}}$ is **coded**.

Corollary: Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{KM+CC}$ and $\mathbb{P} \in \mathcal{C}$ is a class forcing. In the companion model $M_{\mathcal{V}}$, $\mathbb{B}_{\mathbb{P}}$ is an **Ord ^{$M_{\mathcal{V}}$} -cc class-complete Boolean algebra**.

The Class Intermediate Model Theorem

Intermediate Model Theorem: (Solovay)

- If $V \models \text{ZFC}$ and $W \models \text{ZFC}$ is an **intermediate model** between V and its **set-forcing extension** $V[G]$, then W is a **set-forcing extension** of V .
- If $V \models \text{ZF}$ and $V[a] \models \text{ZF}$, with $a \subseteq V$, is an **intermediate model** between V and its **set-forcing extension** $V[G]$, then $V[a]$ is a **set-forcing extension** of V .

Definition: Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{GBC}$. Then $\mathcal{W} = \langle W, \in, \mathcal{C}^* \rangle$ is a **simple extension** of \mathcal{V} if \mathcal{C}^* is **generated** by \mathcal{C} together with a single new class.

- **Forcing extensions** are **simple extensions**.

Definition: Suppose T is a **second-order set theory**.

- The **Intermediate Model Theorem holds for T** if whenever $\mathcal{V} \models T$ and $\mathcal{W} \models T$ is an **intermediate model** between \mathcal{V} and its class-forcing extension $\mathcal{V}[G] \models T$, then \mathcal{W} is a **class-forcing extension** of \mathcal{V} .
- The **simple Intermediate Model Theorem holds for T** if whenever $\mathcal{V} \models T$ and $\mathcal{W} \models T$ is a **simple extension** of \mathcal{V} between \mathcal{V} and its class-forcing extension $\mathcal{V}[G] \models T$, then \mathcal{W} is a **class-forcing extension** of \mathcal{V} .

The Class Intermediate Model Theorem: successes and failures

Theorem:

- (Friedman) The **simple Intermediate Model Theorem for GBC fails**.
- (Hamkins, Reitz) The **simple Intermediate Model Theorem for GBC fails** even for **Ord-cc forcing**.

Theorem: (Antos, Friedman, G.) If $M \models \text{ZFC}^-$ and $M[a] \models \text{ZFC}^-$, with $a \subseteq M$, is an **intermediate model** between M and its **set-forcing extension** $M[G]$, then $M[a]$ is a **set-forcing extension** of M .

- Use the Ord-cc class-complete Boolean completion \mathbb{B}_P .

Corollary: (Antos, Friedman, G.) The **simple Intermediate Model Theorem for KM+CC holds**.

Theorem: (Antos, Friedman, G.) Every model $\mathcal{V} \models \text{KM+CC}$ has a forcing extension $\mathcal{V}[G] \models \text{KM} + \text{CC}$ with a **non-simple intermediate model**. Therefore the **Intermediate Model Theorem for KM+CC fails**.

Question: Does the **simple Intermediate Model Theorem for KM** hold?

Hyperclass forcing

The theory of hyperclass forcing was developed by Antos and Friedman.

Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{KM} + \text{CC}$ and \mathbb{P} is a **hyperclass forcing** of \mathcal{V} .

$G \subseteq \mathbb{P}$ is \mathcal{V} -generic if it meets every dense subhyperclass of \mathbb{P} .

Move to the **companion model** $M_{\mathcal{V}}$, where \mathbb{P} is a **definable class forcing**, and form the **forcing extension** $M_{\mathcal{V}}[G]$.

Suppose $M_{\mathcal{V}}[G] \models \text{ZFC}_I^-$ with κ as the **largest cardinal**.

- \mathbb{P} is **pretame**.
- \mathbb{P} preserves **inaccessibility of κ** .
- e.g, \mathbb{P} is $\text{Ord}^{M_{\mathcal{V}}}$ -cc.

Define $\mathcal{V}[G] = \langle W, \in, \mathcal{C}^* \rangle$:

- $W = V_{\kappa}^{M_{\mathcal{V}}[G]}$
- $\mathcal{C}^* = \{ C \subseteq V_{\kappa}^{M_{\mathcal{V}}[G]} \mid C \in M_{\mathcal{V}}[G] \}$

Useful hyperclass forcing

Suppose $\mathcal{V} = \langle V, \in, \mathcal{C} \rangle \models \text{KM} + \text{CC}$ and $M_{\mathcal{V}}$ is the **companion model** of \mathcal{V} .

Hyperclass Boolean completions

Let $\mathbb{P} \in \mathcal{C}$ be a class forcing.

In $M_{\mathcal{V}}$:

- $\mathbb{B}_{\mathbb{P}}$ is a **definable** $\text{Ord}^{M_{\mathcal{V}}}$ -cc **class-complete Boolean algebra**.
- $\mathbb{B}_{\mathbb{P}}$ is **pretame**.
- \mathbb{P} and $\mathbb{B}_{\mathbb{P}}$ have the **same forcing extensions**.
- The **hyperclass forcing extensions** by $\mathbb{B}_{\mathbb{P}}$ are **precisely forcing extensions** by \mathbb{P} .

Hyperclass forcing and first-order set theory

- (Welch) Assuming large cardinals, models of the form $L[C]$ for a **proper class club C of uncountable cardinals** are characterized as **hyperclass forcing extensions** of a truncated iterate of a mouse with large cardinals.
- (Friedman, G., Müller) Assuming large cardinals, models $L[C_1, \dots, C_n]$ for **specially nested clubs C_i of uncountable cardinals** as **hyperclass forcing extensions** of a truncated iterate of a mouse with stronger large cardinals.
- Uses **class products** of Prikry forcing with **class supports**.