

Ehrenfeucht principles in set theory

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This is joint work with [Gunter Fuchs](#) and [Joel David Hamkins](#) (CUNY).

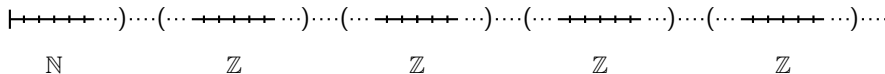
Models of arithmetic

Peano Arithmetic

- The first-order language of arithmetic is $\mathcal{L}_A = \{+, \cdot, <, 0, 1\}$.
- **Peano Arithmetic (PA)** is a collection of statements in \mathcal{L}_A codifying the **fundamental properties of the natural numbers**.
 - ▶ Commutativity and associativity of addition and multiplication, distributive law, ordering is discrete with least element 0, 0 is the additive identity, etc.
 - ▶ **Induction scheme**: for every \mathcal{L}_A -formula $\varphi(x, \vec{y})$,

$$\forall \vec{y}[(\varphi(0, \vec{y}) \wedge (\forall x \varphi(x, \vec{y}) \rightarrow \varphi(x+1, \vec{y}))) \rightarrow \forall x \varphi(x, \vec{y})].$$

- PA proves the **least-number principle**: every definable (with parameters) set has a $<$ -least element.
- The **standard** model of PA is \mathbb{N} .
- A **countable nonstandard** model of PA looks like:

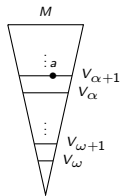


densely many copies of \mathbb{Z}

Models of set theory

Zermelo-Fraenkel set theory (with choice)

- The first-order language of set theory is $\mathcal{L}_S = \{\in\}$.
- **Zermelo-Fraenkel set theory (ZF(C))** is a collection of statements in \mathcal{L}_S codifying the **fundamental properties of sets**.
- A model $M \models \text{ZF}$ is the union of the **von Neumann hierarchy** $M = \bigcup_{\alpha \in \text{ORD}} V_\alpha^M$:
 - ▶ $V_0^M = \emptyset$,
 - ▶ $V_{\alpha+1}^M = P^M(V_\alpha)$ (Powerset of V_α^M),
 - ▶ $V_\lambda^M = \bigcup_{\alpha < \lambda} V_\alpha^M$ for limit ordinals λ .
 - ▶ If $a \in M$, then **$\text{rank}(a) = \alpha$ is least** such that $a \in V_{\alpha+1}^M$.
- **Theorem:** (Lévy-Montague reflection) For every n , there are unboundedly many ordinals α in M such that $V_\alpha^M \prec_{\Sigma_n} M$.



There is an active **exchange of concepts, methods and techniques** between model theory of models of PA and model theory of models of ZF.

Some general model theory

Suppose M is a first-order structure and $P \subseteq M$.

Definition:

- $a \neq b$ in M are **indiscernible** if they have the **same type**:
for every $\varphi(x)$, $M \models \varphi(a) \leftrightarrow \varphi(b)$.
- $a \neq b$ in M are **P -indiscernible** if they have the same type with **parameters from P** .
- Otherwise they are **(P)-discernible**.

It is **common** for models to have indiscernible elements:

- A first-order structure (in a countable language) of **size greater than 2^ω** has indiscernible elements.
- A **non-rigid** first-order structure has indiscernible elements.
- If T is a first-order theory with an infinite model, then there are models of T with sets of indiscernible elements of any cardinality (by compactness).

Ehrenfeucht's lemma in arithmetic

A powerful tool in model theory of models of PA is

Ehrenfeucht's lemma: (Ehrenfeucht, '73) If $a \neq b$ in a model $M \models \text{PA}$ and b is definable from a in M , then a and b are discernible in M .

Here is an application.

Theorem: If $M \models \text{PA}$ is the Skolem closure of a single element $a \in M$, then it has no non-trivial automorphisms.

Proof:

- a is the only element of its type by Ehrenfeucht's lemma.
- An automorphism of M must take a to a and therefore fixes everything. \square

Note: Given any first-order structure, we can ask if Ehrenfeucht's lemma holds there.

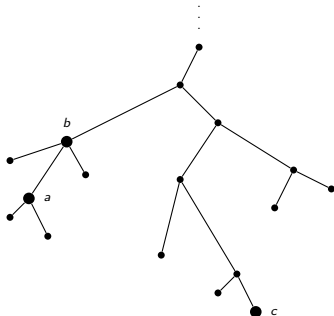
Question: Does Ehrenfeucht's lemma hold for models of ZF?

Proof of Ehrenfeucht's lemma

Proof: Suppose $a \neq b$ and b is definable from a .
Fix a definable function $f(x)$ such that $f(a) = b$.

Case 1: $a < f(a) = b$

- G is the graph whose edges are pairs $\{x, f(x)\}$ such that $x < f(x)$.
- G has an edge between a and b .
- G is loop-free.
- $d(x, y)$ is the length of the shortest path between x and y in G , if connected.
- c is $<$ -least in the connected component of a and b (exists by the least-number principle).
- $d(c, a)$ and $d(c, b)$ differ by 1.
- $d(c, a)$ is even iff $d(c, b)$ is odd.

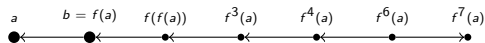


Proof of Ehrenfeucht's lemma (continued)

$f(x)$ is a definable function such that $f(a) = b$.

Case 2: $a > f(a) = b$

- $d(x)$ is the number of times $f(x)$ can be iterated before the values stop **decreasing**.
- $d(x)$ exists by the least-number principle.
- $d(a)$ and $d(b)$ differ by 1.



□

Ordinal definable sets

Suppose $M \models \text{ZF}$ and fix a definable bijection $F : \text{ORD}^{<\omega^M} \xrightarrow[\text{onto}]{1-1} \text{ORD}$.

Definition:

- A set is **ordinal definable** if it is definable with ordinal parameters (using F , we can always assume that there is a **single** ordinal parameter).
- **OD** is the collection of all ordinal definable sets.
- **HOD** is the collection of all hereditarily ordinal definable sets.

Lemma: A set $a \in \text{OD}$ iff M satisfies that a is ordinal definable in some V_α .

Proof:

(\Rightarrow): If $a \in \text{OD}$, then a is defined by the same formula in some V_α (by reflection).

(\Leftarrow): Suppose M satisfies that a is defined by $\ulcorner \varphi(x, \beta) \urcorner$ in V_α .

Note that φ might be **nonstandard**. But

$$\psi(x, \ulcorner \varphi \urcorner, \beta, \alpha) := \exists y \, y = V_\alpha \wedge y \models \ulcorner \varphi(x, \beta) \urcorner$$

defines a in M . \square

Corollary: The collection **OD** is **first-order definable** in M .

A definable well-ordering of OD

Lemma: M has a definable well-ordering $<$ of OD.

Proof: Fix $x, y \in \text{OD}$.

- x is definable in V_α by $\ulcorner \varphi(x, \xi) \urcorner$, where α and $F(\ulcorner \varphi \urcorner, \xi)$ are least.
- y is definable in V_β by $\ulcorner \psi(x, \mu) \urcorner$, where β and $F(\ulcorner \psi \urcorner, \mu)$ are least.

$x < y$ if $\alpha < \beta$ or $\alpha = \beta$ and $F(\ulcorner \varphi \urcorner, \xi) < F(\ulcorner \psi \urcorner, \mu)$. \square

Lemma: HOD is a transitive model of ZFC.

Proof:

- Transitivity is by definition.
- ZF a is direct verification.
- If $a \in \text{OD}$, then the well-ordering of OD restricted to a is ordinal definable.

The axiom: $V = \text{HOD}$

Definition: The axiom $V = \text{HOD}$ states that every set is hereditarily ordinal definable.

Corollary: If $M \models \text{ZF}$, then M is a model of $V = \text{HOD}$ iff M has a **definable well-ordering**.

Proof:

(\Rightarrow): M has a definable well-ordering of ordinal definable sets.

(\Leftarrow): Assume (wlog) order-type of well-ordering is ORD. Each $a \in M$ is the α^{th} -element for some ordinal α . \square

Ehrenfeucht's lemma and ordinal definable sets

Theorem: If $a \neq b$ in a model $M \models \text{ZF}$, a is ordinal definable and b is definable from a in M , then a and b are discernible in M .

“Ehrenfeucht's lemma holds for ordinal definable sets.”

Proof:

- b is ordinal definable.
- Use same argument as for models of PA with the well-ordering $<$ of OD-sets. \square

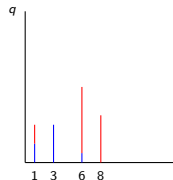
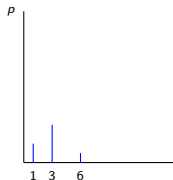
Corollary: Ehrenfeucht's lemma holds for every model of $\text{ZF} + V = \text{HOD}$.

A counterexample to Ehrenfeucht's lemma

Theorem: (Fuchs, G., Hamkins) If $M \models \text{ZFC}$ and $M[c]$ is the Cohen forcing extension, then there are $a \neq b$ in $M[c]$ such that a and b are inter-definable, but a and b are M -indiscernible.

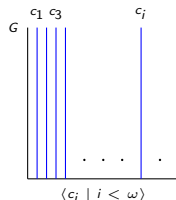
Proof:

- Poset \mathbb{P} to add a Cohen real is isomorphic to finite-support product $\prod_{i < \omega} \mathbb{P}$.
- The poset $\prod_{i < \omega} \mathbb{P}$ adds ω -many Cohen reals:
 - ▶ conditions: finite functions $p : \text{dom}(p) \rightarrow {}^{<\omega}2$ on ω ,
 - ▶ order: $q \leq p$ if $\text{dom}(p) \subseteq \text{dom}(q)$ and for all $n \in \text{dom}(p)$, $q(n)$ extends $p(n)$.



A counterexample to Ehrenfeucht's lemma (continued)

- $G \subseteq \prod_{i < \omega} \mathbb{P}$ is M -generic.
- $C = \{c_i \mid i < \omega\}$.
- $E = \{c_{2i} \mid i < \omega\}$ and $O = \{c_{2i+1} \mid i < \omega\}$.
- $\langle C, E \rangle$ and $\langle C, O \rangle$ are M -indiscernible.



- \dot{C} , \dot{E} , and \dot{O} are \mathbb{P} -names for C , E , and O .
- Suppose $M[G] \models \varphi(\langle C, E \rangle, d)$ with $d \in M$. Then there is $q \Vdash \varphi(\langle \dot{C}, \dot{E} \rangle, \check{d})$ in G .
- If $i \in \text{dom}(q)$ is even/odd, then there is odd/even m_i such that $q(i) \subseteq c_{m_i}$.
- $\pi : \mathbb{P} \rightarrow \mathbb{P}$ is an automorphism such that
 - ▶ π switches even and odd coordinates.
 - ▶ $\pi(i) = m_i$ for all $i \in \text{dom}(q)$.
- $H = \pi " G$ is M -generic for $\prod_{i < \omega} \mathbb{P}$ and $M[G] = M[H]$.
- $q \in H$.
- $\dot{C}_H = C$, $\dot{E}_H = O$, and $\dot{O}_H = E$.
- $M[H] = M[G] \models \varphi(\langle C, O \rangle, d)$. \square

Is Ehrenfeucht's lemma equivalent to $V = \text{HOD}$?

Question: If Ehrenfeucht's lemma holds in $M \models \text{ZF}(C)$, is M a model of $V = \text{HOD}$?

Ehrenfeucht Principles

Parametric generalizations of Ehrenfeucht's lemma.

Definition: (Fuchs, G., Hamkins) Let $M \models \text{ZF}$ and let $A, P, Q \subseteq M$. The principle $\text{EL}(A, P, Q)$ for M asserts that if $a \in A$, $a \neq b$, and b is definable in M from a with parameters from P , then a and b are Q -discernible.

" P -definability from A implies Q -discernibility."

Observation:

- $\text{EL}(M, \emptyset, \emptyset)$ is **Ehrenfeucht's lemma**.
- $\text{EL}(\text{OD}, \emptyset, \emptyset)$ holds for every $M \models \text{ZF}$.
- $\text{EL}(M[c], \emptyset, M)$ **fails** in every Cohen forcing extension $M[c]$.
- $\text{EL}(M, \emptyset, \emptyset)$ implies $\text{EL}(M, P, P)$ for every P .
- $\text{EL}(A, P, Q)$ gets **stronger** if A, P are enlarged or Q is diminished.

Interesting Ehrenfeucht principles

Suppose $M \models \text{ZF}$.

$\text{EL}(M, \text{ORD}, \text{ORD})$

- If $a \neq b$ and b is definable from a with **ordinal parameters**, then a and b are **ORD-discernible**.
- **First-order** expressible (similar argument to definability of OD).

$\text{EL}(M, M, \emptyset)$

- M contains **no indiscernibles**.
- (Leibniz) **Identity of Indiscernibles**: any two objects differ on some property.
- (Enayat) M is **Leibnizian**.
- **Theorem**: (Enayat, '04) There are **uncountable** Leibnizian models of set theory.
- **Theorem**: (Mycielski, '95) A theory T extending ZF has a **Leibnizian model** iff T **proves the Leibniz-Mycielski axiom** (next slide).

$\text{EL}(M, M, \text{ORD})$

- Any $a \neq b$ are **ORD-discernible**.
- **First-order** expressible.

Leibniz-Mycielski axiom

Leibniz-Mycielski axiom (LM): (Mycielski, '95) If $a \neq b$ in $M \models \text{ZF}$, then a and b are discernible in some V_α .

Observation: Suppose $M \models \text{ZF}$.

- LM is equivalent to $\text{EL}(M, M, \text{ORD})$.

Proof: Fix $a \neq b$.

(\Rightarrow) : Suppose LM holds.

- ▶ $V_\alpha \models \varphi(a) \wedge \neg \varphi(b)$.
- ▶ $\psi(x, \alpha) := \exists y y = V_\alpha \wedge y \models \varphi(x)$.
- ▶ $M \models \psi(a, \alpha) \wedge \neg \psi(b, \alpha)$.

(\Leftarrow) : Suppose $\text{EL}(M, M, \text{ORD})$ holds.

- ▶ $M \models \varphi(a, \beta) \wedge \neg \varphi(b, \beta)$ for some ordinal β .
- ▶ $V_\gamma \models \varphi(a, \beta) \wedge \neg \varphi(b, \beta)$ (by reflection).
- ▶ δ codes $\langle \beta, \gamma \rangle$.
- ▶ γ, β are definable in $V_{\delta+1}$ without parameters.
- ▶ a and b are discernible in $V_{\delta+1}$. \square

Leibniz-Mycielski axiom (continued)

Corollary:

- LM is **first-order** expressible.
- LM is equivalent to **LM***: If $a \neq b$ in $M \models \text{ZF}$, then M **satisfies** that a and b are discernible in some V_α .

Observation: $V = \text{HOD} \rightarrow \text{LM}$ over ZF.

Question: Is LM **equivalent** to $V = \text{HOD}$ over $\text{ZF}(C)$?

Leibniz-Mycielski axiom (continued)

Theorem: (Enayat, '04) $M \models \text{ZF} + \text{LM}$ iff M has a definable injection from M into subsets of ordinals

$$F : M \xrightarrow{1-1} {}^{<\text{ORD}}2.$$

Proof:

(\Rightarrow): Suppose LM holds.

- Fix a set a . Let β_a be least above $\text{rank}(a)$ such that $V_{\beta_a} \prec_{\Sigma_2} M$.
- T_a consists of pairs $\langle \ulcorner \varphi \urcorner, \alpha \rangle$ such that $V_\alpha \models \ulcorner \varphi(a) \urcorner$ for some $\alpha < \beta_a$.
- Suppose $a \neq b$ and (wlog) $\text{rank}(a) \leq \text{rank}(b)$.
- $V_\gamma \models \varphi(a) \wedge \neg \varphi(b)$ for some $\gamma < \beta_b$.
- $\langle \ulcorner \neg \varphi \urcorner, \gamma \rangle$ is in T_b but not in T_a .
- $T_a \neq T_b$.
- $F(a) = T_a$ (viewed as subset of ordinals via coding).

(\Leftarrow): Suppose $F : M \xrightarrow{1-1} {}^{<\text{ORD}}2$ and $a \neq b$. Then (wlog) $\alpha \in F(a)$, but $\alpha \notin F(b)$. \square

LM and choice principles

Corollary: If $M \models \text{ZF} + \text{LM}$, then M has a definable linear ordering.

Proof: $<^{\text{ORD}2}$ is linearly ordered lexicographically. \square

The existence of a definable linear ordering is a **weak global choice principle**.

Theorem: (Easton, '64) There are **models of ZFC without a definable linear ordering**.

Proof: Consider the class forcing extension $V[G]$, where a Cohen subset is added to every regular cardinal. \square .

Question: Are there models of **ZFC** having a **definable linear ordering**, but **no definable well-ordering** ($V = \text{HOD}$ fails)?

Question: Is **LM** equivalent over $\text{ZF}(C)$ to the **existence of a definable linear ordering**?

LM and choice principles (continued)

Theorem: (Solovay, '04) **LM does not imply** (even countable) **AC** over ZF.

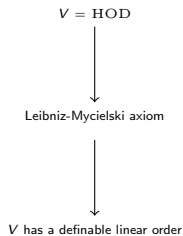
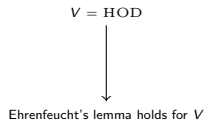
Proof: Similar to Cohen's argument that countable AC fails in $L(\{c_i \mid i < \omega\})$, where c_i are the Cohen reals explicitly added by $\prod_{i < \omega} \mathbb{P}$. But uses Jensen reals.

Work in L .

- **Jensen forcing** \mathbb{Q} :
 - ▶ A **subposet of Sacks forcing** which **has the ccc in L** (constructed using \diamond).
 - ▶ Adds a **unique \mathbb{Q} -generic real over L** .
 - ▶ The **collection of all Jensen reals** in any V is Π^1_2 -definable.
- $\mathbb{Q}^* = \prod_{i < \omega} \mathbb{Q}$ (finite-support).
- **Theorem:** (Kanovei, '14) If $G \subseteq \mathbb{Q}^*$ is L -generic, then the Jensen reals $\{c_i \mid i < \omega\}$ added explicitly by G are the **only Jensen reals in $L[G]$** .
- Consider $L(C)$, where $C = \{c_i \mid i < \omega\}$.
 - ▶ **countable AC fails in $L(C)$** .
 - ▶ C is **definable** in $L(C)$ (as the collection of all Jensen reals).
 - ▶ Every set in $L(C)$ is **ordinal definable** from a unique minimal finite subset of C .
 - ▶ $L(C)$ has a **definable injection into subsets of ordinals**.
 - ▶ $L(C)$ is a **model of LM**. \square

Questions

Question: Suppose $V \models \text{ZFC}$. Which arrows reverse?



Thank you!