Parameter-free schemes in second-order arithmetic

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Second-order arithmetic

Second-order arithmetic has two types of objects: numbers and sets of numbers (reals).

Syntax: Two-sorted logic

- Separate variables and quantifiers for numbers and sets of numbers.
- Convention: lower-case letters for numbers, upper-case letters for sets.
- Notation:
  - Σ_n^0 - first-order Σ_n-formula
  - Σ_n^1 - n-alternations of set quantifiers followed by a first-order formula.

Although first-order logic can be used to formalize second-order arithmetic, that formalization is not natural or easy to work with.

Semantics: A model is $\mathscr{M} = \langle M, +, \times, <, 0, 1, \in, \mathcal{S} \rangle$.

- M is the collection of numbers.
- $\mathcal{S}$ is the collection of sets of numbers: if $A \in \mathcal{S}$, then $A \subseteq M$.

Second-order axioms

- Numbers: Peano Arithmetic PA
- Sets:
  - Extensionality
  - Induction axiom: $\forall X \ ((0 \in X \wedge \forall n(n \in X \to n+1 \in X)) \to \forall n \, n \in X)$

    Applies only to available sets, not second-order formulas.

Weak axiom systems

Arithmetical comprehension ACA_0

Comprehension scheme for first-order formulas: for all n, $\Sigma_n^0\text{-}\mathrm{CA}_0$

if $\varphi(n, A)$ is a first-order formula, then $\{n \mid \varphi(n, A)\}$ is a set.

<small>Note the use of set parameters.</small>

- If $\langle M, +, \times, <, 0, 1 \rangle \models \mathrm{PA}$ and $\mathcal{S}$ consists of definable subsets of M, then

$$\mathcal{M} = \langle M, +, \times, <, 0, 1, \in, \mathcal{S} \rangle \models \mathrm{ACA}_0.$$

- ACA_0 is conservative over PA.

Elementary Transfinite Recursion ATR_0

ACA_0

Transfinite recursion: every first-order recursion on sets along a set well-order has a solution.

- A set well-order is a linear order Γ whose every subset has a minimal element.
- A solution to a recursion is a code of a function $F : \mathrm{dom}(\Gamma) \to \mathcal{S}$.

  <small>A code for F is $\bar{F} = \{\langle n, m \rangle \mid n \in \mathrm{dom}(\Gamma), m \in F(n)\}$.</small>
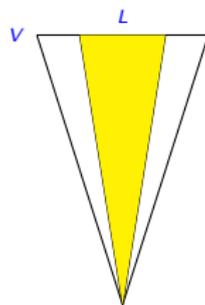
- Applications:
  - ▶ Iterate the Turing jump operation.
  - ▶ Build an internal constructible universe L.

Gödel's L in ATR_0

Gödel's constructible universe L

Suppose $V \models \mathrm{ZF}$.

- $L_0 = \emptyset$
- $L_{\alpha+1}$ is the set of all subsets of L_α definable over L_α.
- $L_\lambda = \bigcup_{\alpha < \lambda} L_\alpha$ for a limit λ.
- $L = \bigcup_{\alpha \in \mathrm{Ord}} L_\alpha$.

Suppose $\mathscr{M} = \langle M, +, \times, <, 0, 1, \in, \mathcal{S} \rangle \models \mathrm{ATR}_0$ and $\Gamma \in \mathcal{S}$ is a well-order.

- $\mathscr{M}$ can construct the L-hierarchy along Γ.
- There is a set coding a sequence of L_Δ for $\Delta \leq \Gamma$ obeying the definition of L.

A model of ATR_0 has its own constructible universe $L^{\mathscr{M}}$!

Comprehension scheme

Increasing the amount of comprehension to more complex second-order assertions produces a hierarchy of second-order set theories.

Σ_n^1-**comprehension** Σ_n^1-CA_0
If $\varphi(n, A)$ is a Σ_n^1-formula, then $\{n \mid \varphi(n, A)\}$ is a set.

- Σ_1^1-CA_0 is stronger than ATR_0.

culminating in:

Full second-order arithmetic $\mathbb{Z}_2$
For all n, Σ_n^1-comprehension.

Definition: **Full parameter-free second-order arithmetic** $\mathbb{Z}_2^{-p}$
Full comprehension without parameters.

Choice scheme

Σ_n^1-choice Σ_n^1-AC

If $\varphi(n, X, A)$ is a Σ_n^1-formula and for every n, there is a set X such that $\varphi(n, X, A)$, then there is a set Y such that for every n, $\varphi(n, Y_n, A)$. $Y_n = \{m \mid \langle n, m \rangle \in Y\}$

Choice scheme AC

For all n, Σ_n^1-choice.

Proposition:

- Over ACA_0, AC implies Z_2.
- The "constructible reals" of a model of Z_2 satisfy $Z_2 + AC$.
- Z_2 is equiconsistent with $Z_2 + AC$.

Definition: **Parameter-free choice scheme** AC^{-p}

Choice scheme without parameters.

Theorem: (Feferman, Lévy) It is consistent that there is a model of Z_2 in which AC^{-p} fails.

Theorem: (Guzicki) It is consistent that there is a model of $Z_2 + AC^{-p}$ in which AC fails.

Collection scheme

Σ_n^1-collection Σ_n^1-Coll
If $\varphi(n, X, A)$ is a Σ_n^1-formula and for every n, there is a set X such that $\varphi(n, X, A)$, then there is a set Y such that for every n, there is m such that $\varphi(n, Y_m, A)$.

Collection scheme Coll
For all n, Σ_n^1-collection.

Proposition:
- Over ACA_0, Coll implies Z_2.
- Over ACA_0, Coll is equivalent to AC.

Definition: **Parameter-free collection scheme Coll^{-p}**
Collection scheme without parameters.

Equiconsistency of Z_2 and Z_2^{-p}

Theorem: (H. Friedman) Z_2 and Z_2^{-p} are equiconsistent.

Proof idea: Suppose $\mathscr{M} = \langle M, +, \times, <, 0, 1, \in, \mathcal{S} \rangle \models Z_2^{-p}$.

- Carry out the L-construction inside $\mathscr{M}$, but use only parameter-free definable well-orders.
- Let $(L^{-p})^{\mathscr{M}}$ be the resulting model.
- The "constructible reals" of $(L^{-p})^{\mathscr{M}}$ satisfy $Z_2 + AC$.

Parameter-free schemes

Questions:

- Is Z_2^{-p} equivalent to Z_2? No
- Does AC^{-p} imply Z_2^{-p} over ACA_0? Open
- Is AC^{-p} equivalent to Coll^{-p} over ACA_0? No
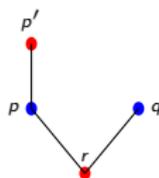
Quick review of forcing

Set-up

- $V \models \mathrm{ZFC}$
- $\mathbb{P} \in V$ is a forcing notion: partial order with largest element **1**

Dense sets and generic filters

$D \subseteq \mathbb{P}$ is dense if for every $p \in \mathbb{P}$, there is $q \in D$ with $q \leq p$.

$G \subseteq \mathbb{P}$ is a filter:

- $\mathbf{1} \in G$.
- (upward closure) If $p \in G$ and $p' \geq p$, then $p' \in G$.
- (compability) If $p, q \in G$, then $r \in G$ such that $r \leq p, q$.

A filter $G \subseteq \mathbb{P}$ is V-generic if it meets every dense set $D \in V$ of $\mathbb{P}$: $D \cap G \neq \emptyset$.

Theorem: V has no V-generic filters for $\mathbb{P}$.

The forcing extension $V[G]$ is constructed from V together with an external V-generic filter G.

Quick review of forcing (continued)

$\mathbb{P}$-**name**: name for an element of $V[G]$

Defined recursively to consist of pairs $\langle \tau, p \rangle$ where $p \in \mathbb{P}$ and τ is a $\mathbb{P}$-name.

Special $\mathbb{P}$-names

- Given $a \in V$, $\check{a} = \{\langle \check{b}, \mathbf{1} \rangle \mid b \in a\}$.
- $\dot{G} = \{\langle \check{p}, p \rangle \mid p \in \mathbb{P}\}$.

Forcing extension $V[G]$

Suppose $G \subseteq \mathbb{P}$ is V-generic.
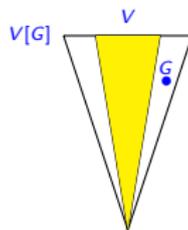
The interpretation of a $\mathbb{P}$-name σ by G:

$\sigma_G = \{\tau_G \mid \langle \tau, p \rangle \in \sigma \text{ and } p \in G\}$. Defined recursively.

The forcing extension $V[G] = \{\sigma_G \mid \sigma \text{ is a } \mathbb{P}\text{-name in } V\}$.

- $V \subseteq V[G]$: $\check{a}_G = a$.
- $G \in V[G]$: $\dot{G}_G = G$.
- $V[G] \models \mathrm{ZFC}$

Forcing relation $p \Vdash \varphi(\sigma)$: whenever G is V-generic and $p \in G$, then $V[G] \models \varphi(\sigma_G)$.

- For a fixed first-order formula $\varphi(x)$, the relation $p \Vdash \varphi(\sigma)$ is definable.
- If $q \leq p$ and $p \Vdash \varphi(\sigma)$, then $q \Vdash \varphi(\sigma)$.
- If $V[G] \models \varphi(\sigma_G)$, then there is $p \in G$ such that $p \Vdash \varphi(\sigma)$.

Cohen forcing

$\mathrm{Add}(\omega, 1)$ - adds a generic real

Elements: binary sequences $p : D \to 2$ with $D \subseteq \omega$ finite.

Order: $q \leq p$ if q extends p.

Suppose $G \subseteq \mathrm{Add}(\omega, 1)$ is V-generic.

- $r = \bigcup G$ is a new real
- $V[G]$ has continuum-many V-generic reals for $\mathrm{Add}(\omega, 1)$.

$$p = \underset{0\,1\,2\,3\,4\,5}{1\,\,1\,0\,\,1}$$

$$q = \underset{0\,1\,2\,3\,4\,5\,6\,7}{1\,0\,1\,0\,\,1\,1\,1}$$

$\mathrm{Add}(\omega, \kappa)$ - adds (at least) κ-many reals

Elements: functions $p : D \to 2$, where D is a finite subset of $\omega \times \kappa$.

Order: $q \leq p$ if q extends p.

Suppose $G \subseteq \mathrm{Add}(\omega, \kappa)$ is V-generic.

$\bigcup G$ gives κ-many new reals.

Sacks forcing

Sacks forcing $\mathbb{S}$ - adds a generic real

Elements: Perfect trees $T \subseteq 2^{<\omega}$: every node has a splitting node above it.

Order: $S \leq T$ if S is a subtree of T.

Suppose $G \subseteq \mathbb{S}$ is V-generic.

- $\bigcap_{T \in G} T = b$ is a branch (real).
- If b is a branch of T, then $T \in G$.
- The generic real b determines G.
- $V[G]$ has continuum-many V-generic reals for $\mathbb{S}$.

Jensen's forcing

Jensen's forcing $\mathbb{J}$ - adds a unique generic real

- $\mathbb{J}$ is a sub-partial order of Sacks forcing $\mathbb{S}$
- constructed using $\diamondsuit$ in L <small>construction is technical</small>
- adds a unique generic Π_2^1-definable singleton real
- used by Jensen to show that it is consistent to have a non-constructible Π_2^1-definable singleton real.

  <small>Every Σ_2^1-definable singleton real is in L by Shoenfield's Absoluteness.</small>

Iterations of forcing notions

Suppose $\mathbb{P}$ is a forcing notion, $G \subseteq \mathbb{P}$ is V-generic, and $\mathbb{Q}$ is a forcing notion in $V[G]$.

V has a $\mathbb{P}$-name $\dot{\mathbb{Q}}$ for $\mathbb{Q}$. Every element of $V[G]$ has a $\mathbb{P}$-name in V.

In V, we define a forcing notion $\mathbb{P} * \dot{\mathbb{Q}}$ such that forcing with $\mathbb{P} * \dot{\mathbb{Q}}$ is the same as forcing with $\mathbb{P}$ followed by forcing with $\mathbb{Q}$.

- Elements: $(p, \dot{q})$ with $p \in \mathbb{P}$ and $p \Vdash \dot{q} \in \dot{\mathbb{Q}}$.
- Order: $(p, \dot{q}) \leq (r, \dot{s})$ if $p \leq r$ and $p \Vdash \dot{q} \leq \dot{s}$.
- *n*-step iterations are defined similarly (infinite iterations can be defined as well).

Example: $\mathbb{S} * \dot{\mathbb{S}}$, where $\dot{\mathbb{S}}$ is the name for the Sacks forcing of the forcing extension.

Sacks forcing of $V[G]$ is different from Sacks forcing of V because $V[G]$ has new perfect trees.

Automorphisms of forcing notions

Suppose $\mathbb{P}$ is a notion of forcing and π is an automorphism of $\mathbb{P}$.

We apply π (recursively) to $\mathbb{P}$-names: $\langle \tau, p \rangle \in \sigma$ if and only if $\langle \pi(\tau), \pi(p) \rangle \in \pi(\sigma)$.

- $\pi(\check{a}) = \check{a}$

The forcing relation respects automorphisms: $p \Vdash \varphi(\sigma)$ if and only if $\pi(p) \Vdash \varphi(\pi(\sigma))$.

If $G \subseteq \mathbb{P}$ is V-generic, then $\pi \text{ " } G$ is V-generic.

Examples

- For any $p, q \in \mathrm{Add}(\omega, 1)$, there is an automorphism π such that p and $\pi(q)$ are compatible.
- Every permutation of κ gives rise to a coordinate-switching automorphism of $\mathrm{Add}(\omega, \kappa)$.
- For any $p, q \in \mathrm{Add}(\omega, \kappa)$, there is a coordinate-switching automorphism π such that p and $\pi(q)$ are compatible.
- Jensen's forcing $\mathbb{J}$ is rigid.

A very bad model of Z_2^{-p}

Theorem: (Kanovei and Lyubetsky) It is consistent that there is a model $\mathcal{M} = \langle M, +, \times, <, 0, 1, \in, \mathcal{S} \rangle \models Z_2^{-p}$ such that $\mathcal{S}$ is not closed under complements.

Proof: Let $G \subseteq \mathrm{Add}(\omega, \omega)$ be V-generic.

- Let $\{a_n \mid n < \omega\}$ be the ω-many generic reals from G.
- Let $\mathcal{S} = P^V(\omega) \cup \{a_n \mid n < \omega\}$.
- $\mathcal{S}$ is not closed under complements.
- Let $\mathcal{M} = \langle \omega, +, \times, <, 0, 1, \in, \mathcal{S} \rangle$.
- Let $\dot{\mathcal{M}}$ be the canonical $\mathrm{Add}(\omega, \omega)$-name for $\mathcal{M}$.
- Fix a second-order formula $\varphi(x)$.
- Suppose, for $n < \omega$, $p \Vdash \dot{\mathcal{M}} \models \varphi(\check{n})$, but $q \Vdash \dot{\mathcal{M}} \models \neg\varphi(\check{n})$.
- Let π be a coordinate-switching automorphism such that there is $r \leq p, \pi(q)$.
- $\pi(\dot{\mathcal{M}}) = \dot{\mathcal{M}}$.
- $\pi(q) \Vdash \pi(\dot{\mathcal{M}}) \models \neg\varphi(\pi(\check{n}))$, and hence $\pi(q) \Vdash \dot{\mathcal{M}} \models \neg\varphi(\check{n})$.
- $r \Vdash \dot{\mathcal{M}} \models \varphi(\check{n})$ ($r \leq p$) and $r \Vdash \dot{\mathcal{M}} \models \neg\varphi(\check{n})$ ($r \leq \pi(q)$). Impossible!
- If some $p \Vdash \dot{\mathcal{M}} \models \varphi(\check{n})$, then all $p \Vdash \dot{\mathcal{M}} \models \varphi(\check{n})$.
- By definability of the forcing relation, $\{n < \omega \mid \mathcal{M} \models \varphi(n)\} \in P^V(\omega) \subseteq \mathcal{S}$.
- $\mathcal{M} \models Z_2^{-p}$. $\square$

A model of $\Sigma_2^1\text{-}CA_0 + Z_2^{-p}$

Question: Are there "nice" models of Z_2^{-p}, but not Z_2?

Theorem: (Kanovei and Lyubetsky) It is consistent that there is a model of $Z_2^{-p} + \Sigma_2^1\text{-}CA_0$ in which $\Sigma_4^1\text{-}CA_0$ fails.

The model is constructed in a forcing extension by a (non-linear) tree iteration of Sacks forcing.

Theorem: (G.) It is consistent that there is a model of $Z_2^{-p} + \Sigma_2^1\text{-}CA_0$ in which $\Sigma_3^1\text{-}CA_0$ fails.

The model is constructed in a forcing extension by a tree iteration of Jensen's forcing.

Finite iterations of Jensen's forcing

Theorem (Abraham) In L, for every $n < \omega$, there is an n-length iteration $\mathbb{J}_n = \mathbb{Q}_0 * \dot{\mathbb{Q}}_1 * \cdots * \dot{\mathbb{Q}}_{n-1}$ such that:

- $\mathbb{Q}_0 = \mathbb{J}$.
- If $G_i \subseteq \mathbb{J}_n \restriction i$ is L-generic, then in $L[G_i]$, $\mathbb{Q} = (\dot{\mathbb{Q}}_i)_{G_i}$ has all properties of Jensen's forcing.
- $\mathbb{J}_n$ adds a unique generic Π^1_2-definable n-length sequence of reals.

Let $\vec{\mathbb{J}} = \langle \mathbb{J}_n \mid n < \omega \rangle$.

The tree iteration of Jensen's forcing along $\omega^{<\omega}$

The tree $\omega^{<\omega}$

$\mathbb{P}(\vec{\mathbb{J}}, \omega^{<\omega})$: tree iteration along $\omega^{<\omega}$

Elements: $p : D_p \to \bigcup_{n<\omega} \mathbb{J}_n$ such that:

- D_p is a finite subtree of $\omega^{<\omega}$,
- for all $s \in D_p$, $p(s) \in \mathbb{J}_{\text{len}(s)}$,
- for $s \subseteq t$ in D_p, $p(s) = p(t) \restriction \text{len}(s)$.

Order: $q \leq p$ if $D_q \supseteq D_p$ and for all $s \in D_p$, $q(s) \leq p(s)$.

Tree iterations of Jensen's forcing

Let T be a tree of height at most ω.

- T is finite.
- $T = \omega_1^{<\omega}$.

$\mathbb{P}(\vec{\mathbb{J}}, T)$: tree iteration along T

Elements: $p : D_p \to \bigcup_{n<\omega} \mathbb{J}_n$ such that:

- D_p is a finite subtree of T,
- for all $s \in D_p$, $p(s) \in \mathbb{J}_{\text{len}(s)}$,
- for $s \subseteq t$ in D_p, $p(s) = p(t) \upharpoonright \text{len}(s)$.

Order: $q \leq p$ if $D_q \supseteq D_p$ and for all $s \in D_p$, $q(s) \leq p(s)$.

Tree iteration of Jensen's forcing (continued)

The tree iteration $\mathbb{P}(\vec{\mathbb{J}}, T)$ adds a tree $\mathcal{T}^G$ (isomorphic to T) such that each node on level n has an L-generic n-length sequence of reals for $\mathbb{J}_n$.

Theorem: (Friedman, G.) Suppose G is L-generic for $\mathbb{P}(\vec{\mathbb{J}}, \omega^{<\omega})$ or $\mathbb{P}(\vec{\mathbb{J}}, \omega_1^{<\omega})$. In $L[G]$:

- The only L-generic n-length sequence of reals for $\mathbb{J}_n$ are those coming from the nodes of $\mathcal{T}^G$.
- The collection of all L-generic n-length sequences of reals for $\mathbb{J}_n$ (any n) is Π_2^1-definable.
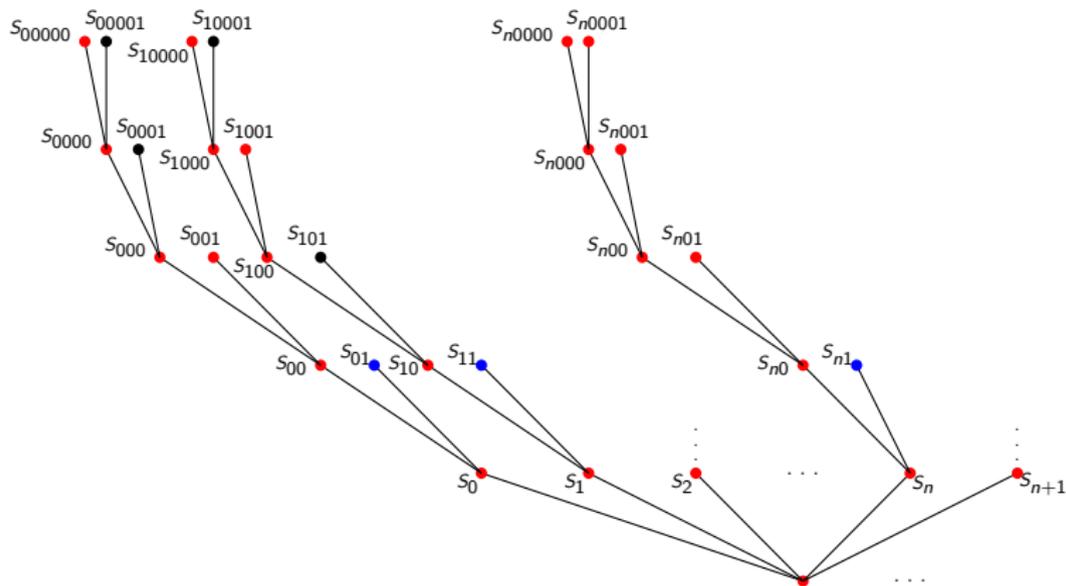
Lemma: Suppose

- $G \subseteq \mathbb{P}(\vec{\mathbb{J}}, \omega^{<\omega})$ (or $\mathbb{P}(\vec{\mathbb{J}}, \omega_1^{<\omega})$) is L-generic.
- T is a finite subtree of $\omega^{<\omega}$ (or countable subtree of $\omega_1^{<\omega}$).
- $G_T = G \upharpoonright \mathbb{P}(\vec{\mathbb{J}}, T)$.

Then

- G_T is L-generic for $\mathbb{P}(\vec{\mathbb{J}}, T)$.
- The only L-generic n-length sequences of reals for $\mathbb{J}_n$ in $L[G_T]$ are those coming from the nodes of T.

Kanovei's tree

Suppose $G \subseteq \mathbb{P}(\vec{\mathbb{J}}, \omega^{<\omega})$ is L-generic. In $L[G]$, define the tree $S \subseteq \mathcal{T}^G$:

- $S_n \in S$ for every $n < \omega$
- $S_{n\vec{0}_m}$ for every $n, m < \omega$

  $\vec{0}_n$ is the sequence of $n \geq 1$-many zeroes

- $S_{n\vec{0}_{m+1}1} \in S$ whenever $S_{n1}(1)(m) = 1$

$S_{01}(1) = 100\ldots$
$S_{11}(1) = 010\ldots$
$S_{n1}(1) = 111\ldots$

A model of $Z_2^{-p} + \Sigma_2^1\text{-CA}_0$ in which $\Sigma_3^1\text{-CA}_0$ fails

Big idea:

- $S_{n1} \notin S$.
- S_{n1} is coded into S.

Let

- $\mathscr{T} = \{ T \subseteq S \mid T \text{ finite} \}$,
- $\mathcal{S} = \{ A \in P^{L[G_T]}(\omega) \mid T \in \mathscr{T} \}$,
- $\mathscr{M} = \langle \omega, +, \times, <, 0, 1, \in, \mathcal{S} \rangle$,
- $\dot{\mathscr{M}}$ is a canonical $\mathbb{P}(\vec{\mathbb{J}}, \omega^{<\omega})$-name for $\mathscr{M}$.

Every permutation f of ω gives rise to an automorphism π_f of $\mathbb{P}(\vec{\mathbb{J}}, \omega^{<\omega})$ which permutes the subtrees $\mathcal{T}_n^G$ (sitting on nodes S_n) of $\mathcal{T}^G$, while preserving the rest of the tree structure.

- $\pi_f(\dot{\mathscr{M}}) = \dot{\mathscr{M}}$
- for any $p, q \in \mathbb{P}(\vec{\mathbb{J}}, \omega^{<\omega})$, there is an automorphism π_f such that p and $\pi_f(q)$ are compatible.

A model of $Z_2^{-p} + \Sigma_2^1\text{-}CA_0$ in which $\Sigma_3^1\text{-}CA_0$ fails (continued)

Theorem (G.): $\mathscr{M} = \langle \omega, +, \times, <, 0, 1, \in, \mathcal{S} \rangle \models \Sigma_2^1\text{-}CA_0 + Z_2^{-p} + \neg\Sigma_3^1\text{-}CA_0$.

Proof:

- $\mathscr{M} \models \Sigma_2^1\text{-}CA_0$ by Shoenfield's Absoluteness.
- $\mathscr{M} \models Z_2^{-p}$ because every parameter-free definable $A \in \mathcal{S}$ is in L by the automorphism argument.
- The collection $\{\vec{r} \in \mathcal{S} \mid \vec{r}$ is an L-generic n-length sequence for $\mathbb{J}_n\} = \{S_s \mid s \in T$ for some $T \in \mathscr{T}\}$ is Π_2^1-definable in $\mathscr{M}$ (uses the construction of the $\mathbb{J}_n$).
- $S_{01}(1) \notin \mathcal{S}$.
- $S_{01}(1)$ is Σ_3^1-definable in $\mathscr{M}$: $m \in S_{01}(1)$ if and only if there are two L-generic $m+1$-length sequences of reals for $\mathbb{J}_{m+1}$ whose first coordinate is S_0. $\square$

Theorem: (G.) Coll^{-p} fails in $\mathscr{M}$.

Proof:

- For every $n < \omega$, $\mathscr{M}$ has an L-generic n-length sequence of reals for $\mathbb{J}_n$.
- If $T \subseteq \mathcal{S}$ is finite, then $L[G_T]$ cannot have a set containing for every $n < \omega$, an L-generic n-length sequence of reals for $\mathbb{J}_n$.

A model of $Z_2^{-p} + \mathrm{Coll}^{-p} + \Sigma_2^1\text{-CA}_0$ in which AC^{-p} and $\Sigma_4^1\text{-CA}_0$ fail

Suppose $G \subseteq \mathbb{P}(\vec{\mathbb{J}}, \omega_1^{<\omega})$ is L-generic. In $L[G]$, define the tree $S \subseteq \mathcal{T}^G$:

- $S_\xi \in S$ for every $\xi < \omega_1$ all level 1 nodes
- $S_{\xi \vec{0}_m}$ for every $\xi < \omega_1$ and $m < \omega$ all left-most branches
- $S_{\xi \vec{0}_{m+1} n}$ for every $\xi < \omega_1$ and $m, k < \omega$. all left-most branch nodes split ω-times
- $S_{\xi \vec{0}_{m+1} \eta} \in S$ for all $\eta < \omega_1$ if and only if $S_{\xi 1}(1)(m) = 1$ left-most branch node splits ω_1-many times

Big idea:

- $S_{\xi 1} \notin S$.
- $S_{\xi 1}$ is coded into S.

Let

- $\mathscr{T} = \{T \subseteq S \mid T \text{ countable}\}$,
- $\mathcal{S} = \{A \in P^{L[G_T]}(\omega) \mid T \in \mathscr{T}\}$,
- $\mathscr{M} = \langle \omega, +, \times, <, 0, 1, \in, \mathcal{S} \rangle$.

Corollary: AC^{-p} and Coll^{-p} are not equivalent over ACA_0 (or even $\Sigma_2^1\text{-CA}_0$).

A model of $Z_2^{-p} + AC^{-p} + \Sigma_2^1\text{-CA}_0$ in which $\Sigma_4^1\text{-CA}_0$ fails

Big idea:

- $S_{\xi n} \notin S$.
- The finite sequences $S_{\xi n} \upharpoonright n$ are coded into S.

Suppose $G \subseteq \mathbb{P}(\vec{\mathbb{J}}, \omega_1^{<\omega})$ is L-generic. In $L[G]$, define the tree $S \subseteq \mathcal{T}^G$:

For $n < \omega$, let p_n be the n-th prime.

- $S_\xi \in S$ for every $\xi < \omega_1$ all level 1 nodes
- $S_{\xi \vec{0}_m}$ for every $\xi < \omega_1$ and $m < \omega$ all left-most branches
- $S_{\xi \vec{0}_{m+1} n}$ for every $\xi < \omega_1$ and $m, k < \omega$. all left-most branch nodes split ω-times
- $S_{\xi \vec{0}_{p_n^m} \eta} \in S$ for all $\eta < \omega_1$ if and only if $S_{\xi n}(1)(m) = 1$ ($m < n$) left-most branch node splits ω_1-many times

Questions

- Can we obtain a model of $Z_2^{-p} + \mathrm{Coll}^{-p} + \Sigma_2^1\text{-}\mathrm{CA}_0$ in which, optimally, $\Sigma_3^1\text{-}\mathrm{CA}_0$ fails?
- Can we obtain a model of $Z_2^{-p} + \mathrm{AC}^{-p} + \Sigma_2^1\text{-}\mathrm{CA}_0$ in which, optimally, $\Sigma_3^1\text{-}\mathrm{CA}_0$ fails?
- Can we obtain a model of $\mathrm{ACA}_0 + \mathrm{AC}^{-p}$ in which Z_2^{-p} fails?
- Given $n < \omega$, can we obtain a model of $Z_2^{-p} + \Sigma_n^1\text{-}\mathrm{CA}_0$ in which $\Sigma_{n+1}^1\text{-}\mathrm{CA}_0$-fails?